

## HENSEL CODES OF SQUARE ROOTS OF $P$ -ADIC NUMBERS

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In this work we are concerned with the calculation of the Hensel codes of square roots of  $p$ -adic numbers, using the fixed point method and this through the calculation of the approached solution of  $f(x) = x^2 - a = 0$  in  $\mathbb{Q}_p$ . We also determine the speed of convergence and the number of iterations.

### 1. INTRODUCTION

The knowledge of the arithmetic and algebraic properties of the  $p$ -adic numbers is useful to the study of their Diophantine properties and the problems of approximations. In this present paper we will see how we can use classical root-finding methods (fixed point) and explore a very interesting application of tools from numerical analysis to number theory. We use this method to calculate the zero of a  $p$ -adic continuous function  $f$  defined on a domain  $D \subset \mathbb{Q}_p$ , where

$$\begin{aligned} f &: \mathbb{Q}_p \rightarrow \mathbb{Q}_p \\ x &\mapsto f(x). \end{aligned}$$

To calculate the square root of a  $p$ -adic number  $a \in \mathbb{Q}_p^*$ , one studies the following problem

$$(1) \quad \begin{cases} f(x) = x^2 - a = 0 \\ a \in \mathbb{Q}_p^*, \quad p - \text{prime number.} \end{cases}$$

Our goal is to calculate the Hensel code of  $\sqrt{a}$ , which means to determine the first numbers of the  $p$ -adic development of the solution of the previous equation,

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and this solution is approached by a sequence of the  $p$ -adic numbers  $(x_n)_n \subset \mathbb{Q}_p^*$  constructed by the fixed point method. We first encountered this idea in [4] where the authors used the numerical methods to find the reciprocal of an integer modulo  $p^n$  (see [4] for more details).

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## 2. PRELIMINARIES

**Definition 2.1.** *Let  $p$  be a prime number. The field of  $p$ -adic numbers  $\mathbb{Q}_p$  is defined as the completion of the field of rational numbers with respect to the  $p$ -adic metric determined by the  $p$ -adic norm. Thus,  $\mathbb{Q}_p$  is obtained from the  $p$ -adic norm in the same way as the real field  $\mathbb{R}$  is obtained from the usual absolute value as the completion of  $\mathbb{Q}$ .*

Here, the function  $|\cdot|_p$  is called the  $p$ -adic norm and is defined by

$$\forall x \in \mathbb{Q}_p : |x|_p = \begin{cases} p^{-v_p(x)} & \text{if } x \neq 0 \\ 0 & \text{if } x = 0, \end{cases}$$

and  $v_p$  is the  $p$ -adic valuation defined by  $v_p(x) = \max\{r \in \mathbb{Z} : p^r \mid x\}$ . The  $p$ -adic distance  $d_p$  is defined by

$$\begin{aligned} d_p & : \mathbb{Q}_p \times \mathbb{Q}_p \rightarrow \mathbb{R}^+ \cup \{0\} \\ (x, y) & \mapsto d_p(x, y) = |x - y|_p. \end{aligned}$$

**Theorem 2.2.** *Every  $p$ -adic number  $a \in \mathbb{Q}_p$  has a unique  $p$ -adic expansion*

$$a = \lambda_n p^n + \lambda_{n+1} p^{n+1} + \dots + \lambda_{-1} p^{-1} + \lambda_0 + \lambda_1 p + \lambda_2 p^2 + \dots = \sum_{k=n}^{\infty} \lambda_k p^k$$

with  $\lambda_k \in \mathbb{Z}$  and  $0 \leq \lambda_k \leq p - 1$  for each  $k \geq n$ .

The short representation of  $a$  is  $\lambda_n \lambda_{n+1} \dots \lambda_{-1} \cdot \lambda_0 \lambda_1 \dots$ , where only the coefficients of the powers of  $p$  are shown. We can use the  $p$ -adic point  $\cdot$  as a device for displaying the sign of  $n$  as follows:

$$\begin{aligned} & \lambda_n \lambda_{n+1} \dots \lambda_{-1} \cdot \lambda_0 \lambda_1 \dots \text{ for } n < 0 \\ & \quad \cdot \lambda_0 \lambda_1 \dots \text{ for } n = 0 \\ & \cdot 00 \dots 0 \lambda_0 \lambda_1 \dots \text{ for } n > 0. \end{aligned}$$

**Definition 2.3.** *A  $p$ -adic number  $a \in \mathbb{Q}_p$  is said to be a  $p$ -adic integer if this canonical development contains only non negative powers of  $p$ . The set of  $p$ -adic*

integers is denoted by  $\mathbb{Z}_p$ . So we have

$$\mathbb{Z}_p = \left\{ \sum_{k=0}^{\infty} \lambda_k p^k, 0 \leq \lambda_k \leq p-1 \right\} = \left\{ a \in \mathbb{Q}_p : |a|_p \leq 1 \right\}.$$

**Definition 2.4.** A  $p$ -adic integer  $a \in \mathbb{Z}_p$  is said to be a  $p$ -adic unit if the first digit  $\lambda_0$  in the  $p$ -adic development is different from zero. The set of  $p$ -adic units is denoted by  $\mathbb{Z}_p^*$ . Hence we have

$$\mathbb{Z}_p^* = \left\{ \sum_{k=0}^{\infty} \lambda_k p^k, \lambda_0 \neq 0 \right\} = \left\{ a \in \mathbb{Q}_p : |a|_p = 1 \right\}.$$

**Proposition 2.5.** Let  $a$  be a  $p$ -adic number. Then  $a$  can be written as the product  $a = p^n \cdot u, n \in \mathbb{Z}, u \in \mathbb{Z}_p^*$ .

**Definition 2.6.** Let  $p$  be a prime number. Then the Hensel code of length  $M$  of any  $p$ -adic number  $a = p^m \cdot u \in \mathbb{Q}_p$  is the pair  $(\text{mant}_a, \text{exp}_a)$ , where the left most  $M$  digits and the value  $m$  of the related  $p$ -adic development are called the mantissa and the exponent, respectively. We use the notation  $H(p, M, a)$  where  $p$  is a prime and  $M$  is the integer which specifies the number of digits of the  $p$ -adic development. One writes

$$H(p, M, a) = (a_m a_{m+1} \dots a_0 a_1 \dots a_t, m),$$

where  $M = |m| + t + 1$ .

See also [1] for more general results concerning the Hensel code.

**Lemma 2.7.** (Hensel's Lemma) Let

$$F(x) = c_0 + c_1 x + c_2 x^2 + \dots + c_n x^n$$

be a polynomial whose coefficients are  $p$ -adic integers. Let

$$F'(x) = c_1 + 2c_2 x + \dots + n c_n x^{n-1}$$

be the derivative of  $F(x)$ . Suppose  $\bar{a}_0$  is a  $p$ -adic integer which satisfies  $F(\bar{a}_0) \equiv 0 \pmod{p}$  and  $F'(\bar{a}_0) \not\equiv 0 \pmod{p}$ . Then there exists a unique  $p$ -adic integer  $a$  such that  $F(a) = 0$  and  $a \equiv \bar{a}_0 \pmod{p}$ .

**Proof.** For the proof of this result we refer the reader to [5]. □

The following theorem makes an important connection between  $p$ -adic numbers and congruences.

**Theorem 2.8.** A polynomial with integer coefficients has a root in  $\mathbb{Z}_p$  if and only if it has an integer root modulo  $p^k$  for every  $k \geq 1$ .

**Proof.** For the proof see [5]. □

A practical consequence of Theorem 2.8 is the following.

**Proposition 2.9.** *A rational integer  $a$  not divisible by  $p$  has a square root in  $\mathbb{Z}_p$  ( $p \neq 2$ ) if and only if  $a$  is a quadratic residue modulo  $p$ .*

**Proof.** Let  $P(x) = x^2 - a$ . Then  $P'(x) = 2x$ . If  $a$  is a quadratic residue, then

$$a \equiv a_0^2 \pmod{p}$$

for some  $a_0 \in \{0, 1, 2, \dots, p-1\}$ . Hence  $P(a_0) \equiv 0 \pmod{p}$ . But

$$P'(a_0) = 2a_0 \not\equiv 0 \pmod{p}$$

automatically since  $(a_0, p) = 1$ , so that the solution in  $\mathbb{Z}_p$  exists by Hensel's lemma. Conversely, if  $a$  is a quadratic nonresidue, by Theorem 2.8 it has no square root in  $\mathbb{Z}_p$ .  $\square$

This can actually be extended to

**Corollary 2.10.** *Let  $p \neq 2$  be a prime. An element  $x \in \mathbb{Q}_p$  is a square if and only if it can be written  $x = p^{2n}y^2$  with  $n \in \mathbb{Z}$  and  $y \in \mathbb{Z}_p^*$  a  $p$ -adic unit.*

**Proof.** For the proof of this result we refer the reader to [3].  $\square$

### 3. MAIN RESULTS

Let  $a \in \mathbb{Q}_p^*$  be a  $p$ -adic number such that

$$|a|_p = p^{-v_p(a)} = p^{-2m}, m \in \mathbb{Z}.$$

If  $(x_n)_n$  is a sequence of  $p$ -adic numbers that converges to a  $p$ -adic number  $\alpha \neq 0$ , then from a certain rank one has

$$|x_n|_p = |\alpha|_p.$$

We also know that if there exists a  $p$ -adic number  $\alpha$  such that  $\alpha^2 = a$ , then  $v_p(a)$  is even and

$$|x_n|_p = |\alpha|_p = p^{-m}.$$

#### 3.1 Fixed point method

To use the fixed point method we study the zeros of the equation  $f(x) = 0$  by studying a related equation  $x = g(x)$ , with the condition that these two formulations are mathematically equivalent. To improve the speed of convergence of the sequence  $(x_n)_n$ , one defines a new sequence that converges more quickly toward the solution of the equation proposed. The conditions that permit the determination of the function  $g(x)$  are:

$$1) g(\sqrt{a}) = \sqrt{a}, g^{(1)}(\sqrt{a}) = \dots = g^{(s-1)}(\sqrt{a}) = 0, g^{(s)}(\sqrt{a}) \neq 0,$$

2) The polynomial  $g(x)$  must not have the square root of  $a$  in its coefficients.

In order to choose  $g(x)$ , we know that if  $\sqrt{a}$  is a root of order  $s$  of  $(g(x) - \sqrt{a})$ , then there is a polynomial  $h(x)$  such that

$$(2) \quad g(x) = \sqrt{a} + (x - \sqrt{a})^s h(x).$$

The conditions that permit the determination of  $h(x)$  are:

- i. The polynomial  $g(x)$  must not have  $\sqrt{a}$  in its coefficients
- ii.  $h(x)$  depends upon the natural number  $s$ .

To determine the formula for  $h(x)$ , it is sufficient to work with the undetermined coefficients and to write the wanted conditions.

Let's consider the following cases.

Case 1:  $s = 1$ . We have

$$g(x) = \sqrt{a} + (x - \sqrt{a})h(x).$$

One chooses  $h(x)$  in order to make the square roots of  $a$  in the coefficients of  $g(x)$  disappear. For this, we put

$$(3) \quad h(x) = \alpha_0.$$

This gives  $\alpha_0 = 1$  and

$$(4) \quad g(x) = x.$$

Case 2:  $s = 2$ . We have

$$(5) \quad g(x) = \sqrt{a} + (x - \sqrt{a})^2 h(x).$$

We put

$$(6) \quad h(x) = \alpha_0 + \alpha_1 x,$$

and get

$$\alpha_0 = -\frac{1}{a^{1/2}}, \quad \alpha_1 = -\frac{1}{2a}.$$

Then we have

$$(7) \quad h(x) = -\frac{1}{2a} (x + 2\sqrt{a}) \text{ and } g(x) = \frac{3}{2}x - \frac{1}{2a}x^3.$$

The sequence associated to  $g(x)$  is defined by

$$(8) \quad \forall n \in \mathbb{N} : x_{n+1} = \frac{3}{2}x_n - \frac{1}{2a}x_n^3.$$

**Theorem 3.1.** *If  $x_{n_0}$  is the square root of  $a$  of order  $r$ , then*

- 1) If  $p \neq 2$ , then  $x_{n+n_0}$  is the square root of  $a$  of order  $2^n r - 2(2^n - 1)m$ .  
 2) If  $p = 2$ , then  $x_{n+n_0}$  is the square root of  $a$  of order  $2^n r - 2(m+1)(2^n - 1)$ .

**Proof.** Let  $(x_n)_n$  the sequence defined by (8). Then

$$(9) \quad \forall n \in \mathbb{N} : x_{n+1}^2 - a = -\frac{1}{4a^2} (4a - x_n^2) (a - x_n^2)^2.$$

We put

$$\Omega(x) = -\frac{1}{4a^2} (4a - x^2).$$

Since

$$(10) \quad |4|_p = \begin{cases} \frac{1}{4} & \text{if } p = 2 \\ 1 & \text{if } p \neq 2, \end{cases}$$

we have

$$\begin{aligned} |\Omega(x_{n_0})|_p &= \left| -\frac{1}{4a^2} (4a - x_{n_0}^2) \right|_p \\ &\leq \left| \frac{1}{4} \right|_p \cdot \left| \frac{1}{a^2} \right|_p \max \{ |4a|_p, |x_{n_0}^2|_p \} \\ &\leq \begin{cases} p^{4m} \cdot \max \{ p^{-2m}, p^{-2m} \} & \text{if } p \neq 2 \\ 2^2 \cdot 2^{4m} \max \{ 2^{-2} \cdot 2^{-2m}, 2^{-2m} \} & \text{if } p = 2 \end{cases} \\ &\leq \begin{cases} p^{2m} & \text{if } p \neq 2 \\ 2^{2m+2} & \text{if } p = 2. \end{cases} \end{aligned}$$

This gives

$$|x_{n_0+1}^2 - a|_p = |\Omega(x_{n_0})|_p \cdot |a - x_{n_0}^2|_p^2,$$

and so we have

$$\begin{cases} |x_{n_0+1}^2 - a|_p \leq p^{2m} \cdot p^{-2r} & \text{if } p \neq 2, \\ |x_{n_0+1}^2 - a|_2 \leq 2^{2m+2} \cdot 2^{-2r} & \text{if } p = 2. \end{cases}$$

Then

$$\begin{cases} x_{n_0+1}^2 - a \equiv 0 \pmod{p^{2r-2m}} & \text{if } p \neq 2, \\ x_{n_0+1}^2 - a \equiv 0 \pmod{2^{2r-2(m+1)}} & \text{if } p = 2. \end{cases}$$

In this manner, we find that if  $p \neq 2$ , then

$$(11) \quad \forall n \in \mathbb{N} : x_{n+n_0}^2 - a \equiv 0 \pmod{p^{\gamma^n}},$$

where the sequence  $(\gamma_n)_n$  is defined by

$$(12) \quad \forall n \in \mathbb{N} : \gamma_n = 2^n r - 2(2^n - 1)m.$$

If  $p = 2$ , then

$$(13) \quad \forall n \in \mathbb{N} : x_{n+n_0}^2 - a \equiv 0 \pmod{2^{\gamma'_n}},$$

where  $(\gamma'_n)_n$  is defined by

$$(14) \quad \forall n \in \mathbb{N} : \gamma'_n = \gamma_n - 2(2^n - 1) = 2^n r - 2(m+1)(2^n - 1).$$

On the other hand, we have

$$(15) \quad \forall n \in \mathbb{N} : x_{n+1} - x_n = \left(-\frac{x_n}{2a}\right) (x_n^2 - a).$$

Since

$$(16) \quad |2|_p = \begin{cases} \frac{1}{2} & \text{if } p = 2 \\ 1 & \text{if } p \neq 2, \end{cases}$$

we have

$$|x_{n+n_0+1} - x_{n+n_0}|_p = \left|\frac{x_{n+n_0}}{2a}\right|_p \cdot |x_{n+n_0}^2 - a|_p.$$

Hence we obtain

$$\begin{cases} |x_{n+n_0+1} - x_{n+n_0}|_p \leq p^{2m} \cdot p^{-m} \cdot p^{-\gamma_n} & \text{if } p \neq 2 \\ |x_{n+n_0+1} - x_{n+n_0}|_2 \leq 2 \cdot 2^{2m} \cdot 2^{-m} \cdot 2^{-\gamma'_n} & \text{if } p = 2, \end{cases}$$

and so

$$\begin{cases} x_{n+n_0+1} - x_{n+n_0} \equiv 0 \pmod{p^{\gamma_n - m}} & \text{if } p \neq 2 \\ x_{n+n_0+1} - x_{n+n_0} \equiv 0 \pmod{2^{\gamma'_n - (m+1)}} & \text{if } p = 2. \end{cases}$$

Therefore, if  $p \neq 2$ , then

$$(17) \quad \forall n \in \mathbb{N} : x_{n+n_0+1} - x_{n+n_0} \equiv 0 \pmod{p^{v_n}},$$

where

$$(18) \quad \forall n \in \mathbb{N} : v_n = \gamma_n - m = 2^n r - (2^{n+1} - 1)m.$$

If  $p = 2$ , then

$$(19) \quad \forall n \in \mathbb{N} : x_{n+n_0+1} - x_{n+n_0} \equiv 0 \pmod{2^{v'_n}},$$

where

$$\forall n \in \mathbb{N} : v'_n = v_n - (2^{n+1} - 1) = 2^n r - (2^{n+1} - 1)(m+1). \quad \square$$

**Conclusion 3.2.**

1. If  $p \neq 2$ , then the following are true.

(a) The speed of convergence of the sequence  $(x_n)_n$  is the order  $v_n$ .

(b) If  $r - 2m > 0$ , then the number of iterations to obtain  $M$  correct digits is

$$(20) \quad n = \left\lceil \frac{\ln \left( \frac{M - m}{r - 2m} \right)}{\ln 2} \right\rceil + 1.$$

(c) With Hensel codes the equation (8) takes the form

$$\begin{aligned} H(p, 2^n r - (2^{n+1} - 1) \cdot m, x) &= H(p, \infty, 3/2) \cdot H(p, 2^{n-1} r - (2^n - 1) \cdot m, x) \\ &\quad - H(p, \infty, 1/2) \cdot \frac{H^3(p, 2^{n-1} r - (2^n - 1) \cdot m, x)}{H^2(p, \infty, x)}. \end{aligned}$$

2. If  $p = 2$ , then the following are true.

(a) The speed of convergence of the sequence  $(x_n)_n$  is the order  $v'_n$ .

(b) If  $r - 2(m + 1) > 0$ , then the necessary number  $n$  of iterations to obtain  $M$  correct digits is

$$(21) \quad n = \left\lceil \frac{\ln \left( \frac{M - (m + 1)}{r - 2(m + 1)} \right)}{\ln 2} \right\rceil + 1.$$

(c) With the Hensel codes the equation (8) takes the form

$$\begin{aligned} H(2, 2^n r - (2^{n+1} - 1) \cdot (m + 1), x) &= H(2, \infty, 3/2) \cdot H(2, 2^{n-1} r - (2^n - 1) \cdot (m + 1), x) \\ &\quad - H(2, \infty, 1/2) \cdot \frac{H^3(2, 2^{n-1} r - (2^n - 1) \cdot (m + 1), x)}{H^2(2, \infty, x)}. \end{aligned}$$

Let's consider for  $p \neq 2$  the sets defined by

$$(22) \quad \begin{cases} S_1 = \left\{ a \in \mathbb{Q}_p : |a|_p = 1 \right\} & \text{if } m = 0 \\ S_2 = \left\{ a \in \mathbb{Q}_p : |a|_p < 1 \right\} & \text{if } m > 0 \\ S_3 = \left\{ a \in \mathbb{Q}_p : |a|_p > 1 \right\} & \text{if } m < 0. \end{cases}$$

We put

$$(23) \quad \forall n \in \mathbb{N} : \begin{cases} v_n^{(1)} = 2^n r & \text{if } m = 0 \\ v_n^{(2)} = 2^n r - (2^{n+1} - 1)m & \text{if } m > 0 \\ v_n^{(3)} = 2^n r - (2^{n+1} - 1)m & \text{if } m < 0. \end{cases}$$



For  $p = 2$ , we consider the sets defined by

$$(24) \quad \begin{cases} B_1 = \{a \in \mathbb{Q}_2 : |a|_2 = 4\} & \text{if } m = -1 \\ B_2 = \{a \in \mathbb{Q}_2 : |a|_2 < 4\} & \text{if } m > -1 \\ B_3 = \{a \in \mathbb{Q}_2 : |a|_2 > 4\} & \text{if } m < -1. \end{cases}$$

We put

$$(25) \quad \forall n \in \mathbb{N} : \begin{cases} v_n^{(1)} = 2^n r & \text{if } m = -1 \\ v_n^{(2)} = 2^n r - (2^{n+1} - 1)(m + 1) & \text{if } m > -1 \\ v_n^{(3)} = 2^n r - (2^{n+1} - 1)(m + 1) & \text{if } m < -1. \end{cases}$$

Then we have the following corollary.

**Corollary 3.3.**

1. If  $p \neq 2$ , then we have the following.

(a) If  $m = 0$ , then we have quadratic convergence for all the  $p$ -adic numbers which belong to the set  $S_1$ .

(b) If  $m < 0$ , then the speed of convergence is faster for all the  $p$ -adic numbers which belong to the set  $S_3$ .

(c) If  $m > 0$ , then the speed of convergence is slower for all the  $p$ -adic numbers which belong to the set  $S_2$ .

2. If  $p = 2$ , then we have the following.

(a) If  $m = -1$ , then one has quadratic convergence for all the 2-adic numbers which belong to  $B_1$ .

(b) If  $m < -1$ , then the speed of convergence is faster for all the 2-adic numbers which belong to the set  $B_3$ .

(c) If  $m > -1$ , then the speed of convergence is slower for all the 2-adic numbers which belong to the set  $B_2$ .

Case 3:  $s = 3$ . We put

$$(26) \quad \begin{cases} g(x) = \sqrt{a} + (x - \sqrt{a})^3 h(x) \\ h(x) = \alpha_0 + \alpha_1 x + \alpha_2 x^2. \end{cases}$$

One finds that

$$(27) \quad \begin{cases} h(x) = \frac{1}{a} + \frac{9}{8a\sqrt{a}} x + \frac{3}{8a^2} x^2 \\ g(x) = \frac{15}{8} x - \frac{5}{4a} x^3 + \frac{3}{8a^2} x^5. \end{cases}$$

The sequence associated to  $g(x)$  is defined by

$$(28) \quad \forall n \in \mathbb{N} : x_{n+1} = \frac{3}{8a^2} x_n^5 - \frac{5}{4a} x_n^3 + \frac{15}{8} x_n.$$

Let  $(x_n)_n$  the sequence defined by (28). Then

$$(29) \quad \forall n \in \mathbb{N} : x_{n+1}^2 - a = (a - x_n^2)^3 \left( -\frac{1}{a^2} + \frac{33}{64a^3} x_n^2 - \frac{9}{64a^4} x_n^4 \right).$$

**Theorem 3.4.** *If  $x_{n_0}$  is the square root of  $a$  of order  $r$ , then the following are true.*

- 1) *If  $p \neq 2$ , then  $x_{n+n_0}$  is the square root of  $a$  of order  $3^n r - 2(3^n - 1)m$ .*
- 2) *If  $p = 2$ , then  $x_{n+n_0}$  is the square root of  $a$  of order  $3^n r - (3^n - 1)(2m + 3)$ .*

**Proof.** For the proof of this theorem, we use the method that we applied in the case where  $s = 2$ .  $\square$

From this theorem, we get that if  $p \neq 2$ , then

$$(30) \quad \forall n \in \mathbb{N} : x_{n+n_0}^2 - a \equiv 0 \pmod{p^{\pi_n}},$$

where  $(\pi_n)_n$  is defined by

$$(31) \quad \forall n \in \mathbb{N} : \pi_n = 3^n r - 2(3^n - 1)m.$$

If  $p = 2$ , then

$$(32) \quad \forall n \in \mathbb{N} : x_{n+n_0}^2 - a \equiv 0 \pmod{2^{\pi'_n}},$$

where  $(\pi'_n)_n$  is given by

$$(33) \quad \forall n \in \mathbb{N} : \pi'_n = 3^n r - (3^n - 1)(2m + 3).$$

On the other hand, we have

$$(34) \quad \forall n \in \mathbb{N} : x_{n+1} - x_n = (a - x_n^2) \left( \frac{7}{8a} x_n - \frac{3}{8a^2} x_n^3 \right).$$

Then if  $p \neq 2$ , we have

$$(35) \quad \forall n \in \mathbb{N} : x_{n+n_0+1} - x_{n+n_0} \equiv 0 \pmod{p^{\Sigma_n}},$$

where  $(\Sigma_n)_n$  is defined by

$$(36) \quad \forall n \in \mathbb{N} : \Sigma_n = 3^n r - (2 \cdot 3^n - 1)m.$$

If  $p = 2$ , then

$$(37) \quad \forall n \in \mathbb{N} : x_{n+n_0+1} - x_{n+n_0} \equiv 0 \pmod{2^{\Sigma'_n}},$$

with

$$(38) \quad \forall n \in \mathbb{N} : \Sigma'_n = 3^n r - ((2 \cdot 3^n - 1)m + 3^{n+1}).$$

**Conclusion 3.5.**

1. If  $p \neq 2$ , then the following are true.

(a) The speed of convergence of the sequence  $(x_n)_n$  is the order  $\Sigma_n$ .

(b) If  $r - 2m > 0$ , then the number  $n$  of necessary iterations to obtain  $M$  correct digits is

$$(39) \quad n = \left\lceil \frac{\ln \left( \frac{M - m}{r - 2m} \right)}{\ln 3} \right\rceil + 1.$$

(c) With Hensel codes the equation (28) takes the form

$$\begin{aligned} & H(p, 3^n r - (2 \cdot 3^n - 1) \cdot m, x) \\ &= H(p, \infty, 3/8) \cdot \frac{H^5(p, 3^{n-1} r - (2 \cdot 3^{n-1} - 1) \cdot m, x)}{H^4(p, \infty, x)} \\ &+ H(p, \infty, -5/4) \cdot \frac{H^3(p, 3^{n-1} r - (2 \cdot 3^{n-1} - 1) \cdot m, x)}{H^2(p, \infty, x)} \\ &+ H(p, \infty, 15/8) \cdot H(p, 3^{n-1} r - (2 \cdot 3^{n-1} - 1) \cdot m, x). \end{aligned}$$

2. If  $p = 2$ , then the following are true.

(a) The speed of convergence of the sequence  $(x_n)_n$  is the order  $\Sigma'_n$ .

(b) If  $r - (2m + 3) > 0$ , then the necessary number of iterations to obtain  $M$  correct digits is

$$(40) \quad n = \left\lceil \frac{\ln \left( \frac{M - m}{r - (2m + 3)} \right)}{\ln 3} \right\rceil + 1.$$

(c) With Hensel codes the equation (28) takes the form

$$\begin{aligned} & H(2, 3^n r - ((2 \cdot 3^n - 1) \cdot m + 3^{n+1}), x) \\ &= H(2, \infty, 3/8) \cdot \frac{H^5(2, 3^{n-1} r - ((2 \cdot 3^{n-1} - 1) \cdot m + 3^n), x)}{H^4(2, \infty, x)} \\ &+ H(2, \infty, -5/4) \cdot \frac{H^3(2, 3^{n-1} r - ((2 \cdot 3^{n-1} - 1) \cdot m + 3^n), x)}{H^2(2, \infty, x)} \\ &+ H(2, \infty, 15/8) \cdot H(2, 3^{n-1} r - ((2 \cdot 3^{n-1} - 1) \cdot m + 3^n), x). \end{aligned}$$

### 3.2 Generalization

Generally we can construct an iterative method that converges to  $\sqrt{a}$  with a higher convergence rate. To accelerate the rate of convergence of the sequence  $(x_n)_n$  as much as one wants, it is necessary to solve the problem of letting

$$(41) \quad g(x) = \sqrt{a} + (x - \sqrt{a})^s h(x, \sqrt{a})$$

and choosing the function  $h(x, \sqrt{a})$  in order to make the square roots of  $a$  in coefficients of a function  $g(x)$  disappear. We take the degree of the function  $h(x)$  equal to  $s - 1$ , giving

$$(42) \quad h(x) = \alpha_0 + \alpha_1 x + \alpha_2 x^2 + \cdots + \alpha_{s-1} x^{s-1} = \sum_{j=0}^{s-1} \alpha_j x^j.$$

Then

$$(43) \quad g(x) = \sqrt{a} + (x - \sqrt{a})^s (\alpha_0 + \alpha_1 x + \alpha_2 x^2 + \cdots + \alpha_{s-1} x^{s-1}) = \sum_{j=0}^{2s-1} c_j(\alpha_i, \sqrt{a}) x^j,$$

where, if  $i \in \{0, \dots, s-1\}$ ,

$$(44) \quad c_j(\alpha_i, \sqrt{a}) = \begin{cases} \sqrt{a} + (-1)^s (\sqrt{a})^s \alpha_0, & \text{if } j = 0 \\ \sum_{i=0}^j \alpha_i \binom{s}{j-i} (-1)^{s-j+i} (\sqrt{a})^{s-j+i}, & \text{if } j \in \{1, \dots, 2s-2\} \\ \alpha_{s-1}, & \text{if } j = 2s-1, \end{cases}$$

and

$$(45) \quad \alpha_i = 0, \quad \forall i > s-1.$$

To generalize the fixed point method it is necessary that the coefficients of the even powers of  $x$  are equal to zero and according to the different calculations that we made, we suppose that this condition is also sufficient until a certain  $s$  sufficiently large, i.e

$$(46) \quad \forall j \in \{0, \dots, s-1\} : \{c_{2k}\}_{k \in \{0, \dots, j\}} = \{0\} \Leftrightarrow \sqrt{a} \notin \{c_{2k+1}\}_{k \in \{0, \dots, j\}}.$$

Therefore, to determine  $(\alpha_i)_{i \in \{0, \dots, s-1\}}$  for any  $s$ , it is sufficient to solve the following linear system

$$(47) \quad \begin{cases} c_0 = \sqrt{a} + (-1)^s (\sqrt{a})^s \alpha_0 = 0 \\ c_2(\alpha_i, \sqrt{a}) = 0 \\ c_4(\alpha_i, \sqrt{a}) = 0 \\ \vdots \\ c_{2s-2}(\alpha_i, \sqrt{a}) = 0. \end{cases}, \quad 0 \leq i \leq s-1$$

We get

$$(48) \quad g(x) = c_1(a)x + c_3(a)x^3 + \dots + c_{2s-1}(a)x^{2s-1}.$$

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