

TWO REPRESENTATION OF REFLEXIVE G -INVERSES AND THEIR IMPLEMENTATION

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In this paper we describe algorithms in the package `ginv` for implementation of two methods for computing reflexive g -inverses. These methods are based on the following general solution of the matrix equations (1) and (2): $G = W_1(QW_1)^{-1}(W_2P)^{-1}W_2$. In the first algorithm we investigate implementation of a general determinantal representation for generalized inverses, which is introduced in [19]. These algorithms are continuation of the analogous algorithms developed in [19], written in the programming language `Mathematica`. The second algorithm describes implementation of a modification of the hyper-power iterative method, introduced in [21].

1. INTRODUCTION

Let \mathbf{C} (resp. \mathbf{R}) be the field of complex (resp. real) numbers and $\mathbf{C}_r^{m \times n}$ (resp. $\mathbf{R}_r^{m \times n}$) be the set of $m \times n$ complex (real) matrices whose rank is r . Conjugate, transposed and conjugate-transposed matrix of A are denoted by \overline{A} , A^T and A^* , respectively. The determinant of a square matrix B is denoted by $|B|$, and $\text{Tr}(A)$ denotes the trace of A .

For a given $m \times n$ matrix A over \mathbf{C} , let $\alpha = \{\alpha_1, \dots, \alpha_r\}$ and $\beta = \{\beta_1, \dots, \beta_r\}$ be subsets of $\{1, \dots, m\}$ and $\{1, \dots, n\}$, respectively. Then $|A_\beta^\alpha|$ denotes the minor of A determined by the rows indexed by α and the columns indexed by β .

We use the following notation from [10]. For $1 \leq k \leq n$, denote the collection of strictly increasing sequences of k integers chosen from $\{1, \dots, n\}$, by

$$\mathcal{Q}_{k,n} = \{\alpha : \alpha = (\alpha_1, \dots, \alpha_k), 1 \leq \alpha_1 < \dots < \alpha_k \leq n\}.$$

Let $\mathcal{N} = \mathcal{Q}_{r,m} \times \mathcal{Q}_{r,n}$. For fixed $\alpha \in \mathcal{Q}_{k,m}$, $\beta \in \mathcal{Q}_{k,n}$, $1 \leq k \leq r$, let

$$\mathcal{I}(\alpha) = \{I : I \in \mathcal{Q}_{r,m}, I \supseteq \alpha\}, \quad \mathcal{J}(\beta) = \{J : J \in \mathcal{Q}_{r,n}, J \supseteq \beta\},$$

$$\mathcal{N}(\alpha, \beta) = \mathcal{I}(\alpha) \times \mathcal{J}(\beta).$$

If A is a square matrix, then the coefficient of a_{ij} in the LAPLACE's expansion of $|A|$ is denoted by $\frac{\partial}{\partial a_{ij}}|A|$.

For any matrix $A \in \mathbf{C}^{m \times n}$ the MOORE-PENROSE inverse of A is the unique matrix, denoted by A^\dagger , satisfying the following PENROSE's equations in X :

$$(1) \quad AXA = A, \quad (2) \quad XAX = X, \quad (3) \quad (AX)^* = AX, \quad (4) \quad (XA)^* = XA$$

and if $m = n$, also

$$(5) \quad AX = XA.$$

For a sequence \mathcal{S} of $\{1, 2, 3, 4, 5\}$, the set of matrices obeying the conditions contained in \mathcal{S} is denoted by $A\{\mathcal{S}\}$. A matrix from $A\{\mathcal{S}\}$ is called an \mathcal{S} -inverse of A and denoted by $A^{(\mathcal{S})}$. In the case $m = n$, the group inverse of A , denoted by $A^\#$, is the unique $\{1, 2, 5\}$ inverse of A .

Main properties of the weighted MOORE-PENROSE inverse are investigated in [5], [13]. By $A_{M,N}^\dagger$ we denote the unique solution of the equations (1), (2) and the following equations:

$$(6) \quad (MAX)^* = MAX \quad (7) \quad (NXA)^* = NXA.$$

The methods implemented in this paper are based on the general representations of different classes of pseudoinverses, investigated in [5], [7], [13], [15], [18].

The paper is organized as follows. The second section contains implementation of the determinantal representation of generalized inverses, considered in [4], [19], [20], [22], so called the *general determinantal representation*. This implementation represents a continuation of the paper [19], where a general determinantal representation for the class of $\{1, 2\}$ inverses is introduced. Also, in [19] are developed algorithms in the programming language **C** for implementation of the general determinantal representation.

In the third section we describe implementation in **MATHEMATICA** of an iterative method for computing $\{1, 2\}$ inverses. This method is introduced in [21], and it is based on a generalization of the hyper-power method.

Several illustrative examples are given in the last section.

In this way, we obtain an extension of the programming system **MATHEMATICA**, by means of the implemented functions for computing the rank and index of a given matrix, and by means of the functions for computing the following classes of pseudoinverses: MOORE-PENROSE, weighted MOORE-PENROSE inverse, group inverse, $\{1, 2, 3\}$, $\{1, 2, 4\}$, $\{1, 2\}$ inverses, left/right inverses, RADIC's and STOJAKOVIĆ's (JOSHI's) generalized inverses. It is well known that in **MATHEMATICA** is available only the function *PseudoInverse* for computing the MOORE-PENROSE inverse [25], [26].

2. IMPLEMENTATION OF GENERAL DETERMINANTAL REPRESENTATION

For the sake of completeness, we restate here the general determinantal representation from the articles [19], [20]. This representation can be obtained from $G = W_1(QW_1)^{-1}(W_2P)^{-1}W_2$, where $A = PQ$ is an arbitrary full-rank factorization of A .

Proposition 2.1. *Let $A \in \mathbf{C}_r^{m \times n}$, and $A = PQ$ be its full-rank factorization. Then an arbitrary $\{1, 2\}$ inverse $G = (g_{ij})$ of A can be represented by the following determinantal representation:*

$$g_{ij} = \frac{\sum_{(\alpha, \beta) \in \mathcal{N}(j, i)} |(W_1 W_2)^T{}_{\beta}^{\alpha}| \frac{\partial}{\partial a_{ji}} |A_{\beta}^{\alpha}|}{\sum_{(\gamma, \delta) \in \mathcal{N}} |(W_1 W_2)^T{}_{\delta}^{\gamma}| |A_{\delta}^{\gamma}|}, \quad \left(\begin{array}{l} 1 \leq i \leq n \\ 1 \leq j \leq m \end{array} \right),$$

where $W_1 \in \mathbf{C}^{n \times r}$ and $W_2 \in \mathbf{C}^{r \times m}$ satisfy the conditions (1.1).

In [19] we introduce notions of the generalized determinant and the *general determinantal representation of different orders*. For the sake of clarity we introduce several notations.

The set $\mathcal{Q}_{t, m} \times \mathcal{Q}_{t, n}$, $t \leq r = \text{rank}(A)$ is denoted by $\mathcal{N}(t)$. For as given $m \times n$ complex matrix R , the *generalized determinant of the order t* , denoted by $\text{DET}_{(R, t)}(A)$, can be expressed in this way (see [19]):

$$(2.1) \quad \text{DET}_{(R, t)}(A) = \sum_{(\gamma, \delta) \in \mathcal{N}(t)} |\overline{R}_{\delta}^{\gamma}| |A_{\delta}^{\gamma}|.$$

Also, we introduce the following notation:

$$\mathcal{I}(\alpha, t) = \{I : I \in \mathcal{Q}_{t, m}, I \supseteq \alpha\}, \quad \mathcal{J}(\beta, t) = \{J : J \in \mathcal{Q}_{t, n}, J \supseteq \beta\},$$

$$\mathcal{N}(\alpha, \beta, t) = \mathcal{I}(\alpha, t) \times \mathcal{J}(\beta, t), \quad t \leq r.$$

Then the *general determinantal representation of the order t* for $A \in \mathbf{C}_r^{m \times n}$, can be written as follows (see [19]):

$$(2.2) \quad g_{ij} = a_{ij}^{(\dagger, R, t)} = \frac{\sum_{(\alpha, \beta) \in \mathcal{N}(j, i, t)} |\overline{R}_{\beta}^{\alpha}| \frac{\partial}{\partial a_{ji}} |A_{\beta}^{\alpha}|}{\text{DET}_{(R, t)}(A)}, \quad \left(\begin{array}{l} 1 \leq i \leq n \\ 1 \leq j \leq m \end{array} \right).$$

In the case $t = r$, we obtain well known result from [4]. Definitions of the generalized determinant and the general determinantal representation of different orders are useful in the case when the rank of a given matrix A is unknown. Then we start the computation using (2.2) with $t = \min\{m, n\}$. Then we decrease values for t , until a nonzero value of t satisfying $\text{DET}_{(R, t)}(A) \neq 0$ is reached.

The general determinantal representation contains all known determinantal representations of generalized inverses, introduced in the papers [3–6], [8], [12–13], [16–17], [23–24].

A few connections between the general determinantal representation and the corresponding results from [10], [14] are investigated in [22].

Now, we describe implementation of the general determinantal representation of different orders, in the package **MATHEMATICA**. We begin by several auxiliary procedures.

2.1. By means of the following routine can be detected square matrices.

```
SquareMatrixQ[a_]:=Length[a]==Length[a[[1]]] /; MatrixQ[a]
```

2.2. The rank of any given matrix A is equal to the number of nonzero elements in the reduced row echelon form of A . The result of the expression $zeros[u]$ is 0 if the vector u is identical to the corresponding zero vector, and 1 otherwise.

```
zeros[u_]:=
  Block[{v=u,n,i=1,lg=0},
    n=Length[v];
    While[i<=n,
      If[v[[i]] != 0, lg=1];
      i++;
    ];
    lg
  ];
```

The function $rank[a]$ is a counter of all nonzero rows contained in the reduced row echelon form of A .

```
rank[a_]:=
  Block[{b=a,i,m,n,r,c},
    {m,n}=Dimensions[b];
    b=RowReduce[b];
    r=Sum[zeros[b[[i]]], {i,m}]
  ]; MatrixQ[a]
```

The index of a square matrix A is defined as the first integer k satisfying $rank(A^{k+1})=rankA^k$.

```
Index[a_]:=
  Block[{b=a,c=IdentityMatrix[Length[a]],d=a,k=0},
    While[Rank[c]!=Rank[d],
      d=d.b; c=c.b; k+=1
    ];
    k
  ] /; SquareMatrixQ[a]
```

2.3. The generalized determinant of the order t , defined by $\text{DET}_{(R,t)}(A)$ in (2.1), can be computed by means of the following procedure $GDetR$. The formal parameters a and r denote the matrices A and R , respectively, and the parameter t denotes the size of the selected minors.

```
GDetR[a_,r_, t_Integer]:=
  Block[{b=a, ra=r, f, s, k, l, ma,mc},
    ma=Minors[b,t]; mc=Minors[ra,t];
    {f,s}=Dimensions[ma];
    Sum[Conjugate[mc[[k,l]]] ma[[k,l]], {k,f}, {l,s}]
  ]/; MatrixQ[a] && MatrixQ[r] &&
    Dimensions[a]==Dimensions[r] && Rank[a]==Rank[r]
```

2.4. In order to implement the general determinantal representation, firstly we develop two useful functions. The first function generates the submatrix of a given matrix A , obtained by deleting its i -th row and j -th column.

```
MatrixComp[a_, i_Integer, j_Integer]:=
  Block[{b=a},
    b=Drop[b,{i,i}]
    b=Transpose[Drop[Transpose[b],{j,j}]];
  ]/; MatrixQ[a]
```

In the second function we generate the submatrix of A determined by the rows p_1, \dots, p_t and columns q_1, \dots, q_t .

```
Minor[a_, p_List, q_List, t_Integer]:=
  Block[{b=a,i,j, c},
    c=IdentityMatrix[r];
    For[i=1, i<=t, i++,
      For[j=1, j<=t, j++,
        c[[i,j]]=b[[p[[i]],q[[j]]]]
      ]
    ];
  ]/; MatrixQ[a]
```

2.5. Using an algorithm from [9], the set of all combinations of the order t of the set $\{1, \dots, n\}$ can be implemented by the following code:

```
While[j>=1,
  If[j>=1,
    For[i=t, i>=j, i--,
      p[[i]]=p[[j]]+i-j+1; p1[[i]]=p[[i]]
    ]
  ]
];
```

2.6. Finally, in the procedure $RINVERSE$ we implement the general determinantal representation of the order $t \leq r$, given by (2.2). The formal parameters a and r represent the input matrices A and R , respectively. Initial value for the order t of selected minors is $t = \min\{m, n\}$. In the while cycle the value of t is decreased until the conditions $\text{DET}_{(R,t)}(A) \neq 0$ is satisfied.

```

RInverse[a_,r_]:=
Block[{b=a,ra=r,t,p,q,m,n,w,v,i,j,k,j1,p1,q1,
      pr,pr1,awv,mr,mrr,mc,s,inv,sw,am},
  inv=Transpose[b]; {m,n}=Dimensions[b];
  t=Min[m,n]; d=GDetR[b,ra,t];
  While[d==0, d=GDetR[b,ra,t]; t-- ];
  p=q=Range[t]; p1=q1=q;
  For[v=1, v<=n, v++,
    For[w=1, w<=m, w++,
      s=0;
      If[t==m, j=1, j=m];
      While[j>=1,
        If[t==n, j1=1, j1=n];
        While[j1>=1,
          pr=pr1=1;
          While[pr<=t && p[[pr]]!=w, pr++];
          While[pr1<=t && q[[pr1]]!=v, pr1++];
          If[pr<=t && pr1<=t,
            mr=Minor[b,p,q,t];
            mrr=Minor[ra,p,q,t];
            mc=Conjugate[Det[mrr]];
            am=Det[MatrixComp[mr,pr,pr1]];
            awv=(-1)^(pr+pr1) am mc,
            awv=0
          ];
          s+=awv;
          If[q[[t]]==n, j1--, j1=t ];
          If[j1>=1,
            For[i=t, i>=j1, i--,
              q[[i]]=q[[j1]]+i-j1+1;
              q1[[i]]=q[[i]]
            ]
          ];
          q1=q=Range[t];
          If[p[[t]]==m, j--, j=t ];
          If[j>=1,
            For[i=t, i>=j, i--,
              p[[i]]=p[[j]]+i-j+1; p1[[i]]=p[[i]]
            ]
          ];
          inv[[v,w]]=s/d
          p=q=Range[t]; p1=q1=q
        ] ];
  inv
]; MatrixQ[a]

```

Remark 2.2. *Described algorithms in MATHEMATICA are simpler and more efficient with respect to the corresponding in [19], written in C.*

Computation of generalized inverses by means of the general determinantal representation is a direct method, and does not use the Gaussian elimination.

3. MODIFICATION OF THE HYPER-POWER METHOD

The hyper-power iterative method was originally devised by ALTMAN [2] for inverting of a nonsingular bounded operator in a BANACH space. In [11] the convergence of the same method is proved under the condition which is weaker than the one assumed in [2], and some better error estimates are derived. ZLOBEC in [30] defined two hyper-power iterative methods of an arbitrary high order $q \geq 2$.

In the paper [21] we adapt the hyper-power method to be valid for computing all of the reflexive g -inverses.

Proposition 3.1. (see [21]) *Let $\text{rank}(A) = r \geq 2$, and the matrices $W_1 \in \mathbf{C}^{n \times r}$, $W_2 \in \mathbf{C}^{r \times m}$ satisfy conditions (1.1). If $q \geq 2$ is an integer, then both of the following two iterative methods:*

$$Y_0 = Y'_0 = \alpha(W_2AW_1)^*, \quad 0 < \alpha \leq \frac{2}{\text{Tr}((W_2AW_1)^*W_2AW_1)},$$

$$\begin{cases} T_k = I_r - Y_kW_2AW_1, \\ Y_{k+1} = (I_r + T_k + \dots + T_k^{q-1})Y_k, \\ X_{k+1} = W_1Y_{k+1}W_2 \end{cases}, \quad \begin{cases} T'_k = I_r - W_2AW_1Y'_k, \\ Y'_{k+1} = Y'_k(I_r + T'_k + \dots + T_k^{q-1}), \\ X'_{k+1} = W_1Y'_{k+1}W_2 \quad k = 0, 1, \dots \end{cases}$$

generate the class of the reflexive g -inverses of A .

Under the suitable conditions, we get iterative methods for computing $\{1, 2, 3\}$ or $\{1, 2, 4\}$ inverses, the MOORE-PENROSE inverse, weighted MOORE-PENROSE inverse or the group inverse of A (see [21]).

3.1. Implementation of the modified hyper-power method is given in the following. In order to compute the value $\alpha = \frac{2}{\text{Tr}((W_2AW_1)^*W_2AW_1)}$, we need a function for computing the trace of a square matrix. This function is not built-in in MATHEMATICA. For this purpose we can use the following one-liner idea from [1]:

```
trace[mat_?MatrixQ]:=
  Plus @@(IdentityMatrix[Length[mat]] mat // Flatten)
```

We recommend the following routine:

```
trace[a_]:=
  Block[{b=a, i},
    Sum[b[[i,i]], {i,Length[b]}]
  ]/; SquareMatrixQ[a]
```

3.2. Now, we give the following implementation of the modified hyper-power method. In the following procedure the parameters a , $w1$, $w2$ represent the matrices A , W_1 , W_2 , respectively. The parameter q denotes the order of the hyper-power expansion, and $numit$ denotes the number of iterations.

```

HyperPower[a_,w1_,w2_,q_,numit_]:=
Block[{tk,tk1,b=a,wa=w1,wb=w2,e,alpha,x,y,c=wb.b.wa, ra,s,i,k=1},
  ra=rank[b];
  alpha=2/trace[Conjugate[Transpose[c]].c];
  y=alpha Conjugate[Transpose[c]];
  e=IdentityMatrix[ra];
  While[k<numit,
    tk1=tk=e-y.c; s=e;
    Do[s+=tk; tk=tk1.tk,{i,q-1}];
    y=s.y; x=wa.y.wb; k+=1
  ];
  x
]

```

4. EXAMPLES

Example 4.1. Consider the test matrix S_5 from [27], in the case $a = 1$, i.e. $S_5 =$

$\begin{pmatrix} 2 & 1 & 1 & 1 & 2 \\ 1 & 0 & 1 & 1 & 1 \\ 1 & 1 & 2 & 1 & 1 \\ 1 & 1 & 1 & 0 & 1 \\ 2 & 1 & 1 & 1 & 2 \end{pmatrix}$. Its full-rank factorization is, for example:

$$P = \begin{pmatrix} 2 & 1 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 2 & 1 & 1 \\ 1 & 1 & 1 & 0 \\ 2 & 1 & 1 & 1 \end{pmatrix}, \quad Q = \begin{pmatrix} 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \end{pmatrix}.$$

For the matrices W_1 and W_2 we can select, for example

$$W_1 = \begin{pmatrix} 1 & 2 & 5 & 3 \\ -2 & 4 & 0 & 3 \\ 2 & 1 & 0 & -2 \\ 0 & 5 & 0 & 1 \\ 7 & 2 & -3 & 2 \end{pmatrix}, \quad W_2 = \begin{pmatrix} 2 & -2 & 1 & 1 & -5 \\ 3 & 0 & 1 & 4 & 0 \\ 0 & 2 & 1 & 3 & 4 \\ 7 & 1 & 1 & 9 & -3 \end{pmatrix}.$$

$RInverse[S_5, Transpose[W_1.W_2]]$ gives

$$S_5^{(1,2)} = \begin{pmatrix} \frac{28759}{10220} & -\frac{113}{140} & -\frac{1}{28} & \frac{151}{20} & -\frac{61317}{10220} \\ \frac{472}{73} & -2 & 1 & -1 & -\frac{399}{73} \\ -\frac{300}{73} & 1 & 0 & 1 & \frac{227}{73} \\ \frac{186}{73} & -1 & 1 & -2 & -\frac{113}{73} \\ -\frac{46539}{10220} & \frac{253}{140} & -\frac{27}{28} & -\frac{131}{0} & \frac{79097}{10220} \end{pmatrix}.$$

$RInverse[S_5, P.Transpose [W_1]]$ gives

$$S_5^{(1,2,3)} = \begin{pmatrix} \frac{223}{140} & -\frac{113}{140} & -\frac{1}{28} & \frac{151}{20} & -\frac{123}{140} \\ \frac{1}{2} & -2 & 1 & -1 & \frac{1}{2} \\ -\frac{1}{2} & 1 & 0 & 1 & -\frac{1}{2} \\ \frac{1}{2} & -1 & 1 & -2 & -\frac{1}{2} \\ -\frac{223}{140} & \frac{253}{140} & -\frac{27}{28} & -\frac{131}{20} & \frac{223}{140} \end{pmatrix}.$$

$RInverse[S_5, Transpose [W_2].Q]$ gives

$$S_5^{(1,2,4)} = \begin{pmatrix} -\frac{127}{146} & \frac{1}{2} & -\frac{1}{2} & \frac{1}{2} & \frac{127}{146} \\ \frac{472}{73} & -2 & 1 & -1 & -\frac{399}{73} \\ -\frac{300}{73} & 1 & 0 & 1 & \frac{227}{73} \\ \frac{186}{73} & -1 & 1 & -2 & -\frac{113}{73} \\ -\frac{127}{146} & \frac{1}{2} & -\frac{1}{2} & \frac{1}{2} & \frac{127}{14} \end{pmatrix}.$$

The value of the expression $RInverse[S_5, S_5]$ is

$$S_5^\dagger = \begin{pmatrix} 0 & \frac{1}{2} & -\frac{1}{2} & \frac{1}{2} & 0 \\ \frac{1}{2} & -2 & 1 & -1 & \frac{1}{2} \\ -\frac{1}{2} & 1 & 0 & 1 & -\frac{1}{2} \\ \frac{1}{2} & -1 & 1 & -2 & \frac{1}{2} \\ 0 & \frac{1}{2} & -\frac{1}{2} & \frac{1}{2} & 0 \end{pmatrix},$$

which is well-known result in [27].

Example 4.2. Consider the matrix $A = \begin{pmatrix} 1 & 0 \\ -1 & 0 \\ 0 & 1 \end{pmatrix}$. Its full-rank factorization is $P = A$, $Q = I_2$. If we select $W_1 = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$, and $W_2 = \begin{pmatrix} 1 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix}$, application of the modified hyper-power method of the order 2 leads to:

$$\begin{aligned}
X_1 &= \begin{pmatrix} 0 & -\frac{4}{49} & 0 \\ \frac{4}{49} & \frac{4}{49} & \frac{4}{49} \end{pmatrix}, & X_2 &= \begin{pmatrix} 0 & -\frac{376}{2401} & 0 \\ \frac{376}{2401} & \frac{376}{2401} & \frac{376}{2401} \end{pmatrix}, \\
X_3 &= \begin{pmatrix} 0 & -\frac{1664176}{5764801} & 0 \\ \frac{1664176}{5764801} & \frac{1664176}{5764801} & \frac{1664176}{5764801} \end{pmatrix}, \\
X_4 &= \begin{pmatrix} 0 & -\frac{16417805178976}{33232930569601} & 0 \\ \frac{16417805178976}{33232930569601} & \frac{16417805178976}{33232930569601} & \frac{16417805178976}{33232930569601} \end{pmatrix}, \\
X_5 &= \begin{pmatrix} 0 & -x_5 & 0 \\ x_5 & x_5 & x_5 \end{pmatrix}, & x_5 &= \frac{821679232341479087467408576}{1104427674243920646305299201}.
\end{aligned}$$

We have obtained sequence converging to $X = \begin{pmatrix} 0 & -1 & 0 \\ 1 & 1 & 1 \end{pmatrix} \in A\{1, 2\}$. The matrix A is of full column rank, so that $A^{(1,2)} = A^{(1,2,4)}$ [29], and consequently $X \in A\{1, 2, 4\}$.

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