

## A CONTRIBUTION TO THE STUDY OF REAL GRAPH POLYNOMIALS

*Ivan Gutman*

Let  $G$  be a graph and  $C$  its circuit. Let  $\alpha$  stand for the matching polynomial. Let  $t$  be a real number, such that  $|t| \leq 1$ . We demonstrate that for certain classes of compound graphs the polynomial  $\beta_t(G, C, x) = \alpha(G, x) - 2t\alpha(G - C, x)$  is real, i.e. that all its zeros are real.

### 1. INTRODUCTION

Let  $G$  be a graph containing  $n$  vertices and  $m$  edges. For  $k > 1$  let  $m(G, k)$  denote the number of ways in which  $k$  independent edges can be selected in  $G$ . In addition to this, let  $m(G, 0) = 1$  and  $m(G, 1) = m$ . Then the matching polynomial of  $G$  is defined as

$$\alpha(G) = \alpha(G, x) = \sum_{k \geq 0} (-1)^k m(G, k) x^{n-2k}.$$

The theory of the matching polynomials is nowadays well elaborated [2]. In particular, it is known that  $\alpha(G)$  is a real graph polynomial [4], i.e. that all its zeros are real-valued numbers.

Consider now another polynomial defined as

$$\beta(G, C) = \beta(G, C, x) = \alpha(G, x) - 2\alpha(G - C, x)$$

where  $C$  is a circuit of the graph  $G$  and  $G - C$  is the subgraph obtained by deleting all the vertices of  $C$  from  $G$ ; if  $C$  is a Hamiltonian circuit of  $G$ , then  $\alpha(G - C) \equiv 1$ .

The zeros of the polynomial  $\beta$  play a distinguished role in AIHARA's theory of cyclic conjugation [1], [12]. It has been recently conjectured [8] that the polynomial  $\beta(G, C)$  is real for all cyclic graphs  $G$  and all circuits  $C$  contained in them. The conjecture was supported by proving its validity for various classes of graphs [7], [8], [11], [12].

In the case of unicyclic graph  $\beta(G, C)$  coincides with the characteristic polynomial. Hence we have the following elementary result.

**Observation 1.** *If  $G$  is a unicyclic graph, then  $\beta(G, C)$  is a real polynomial.*

The above observation may serve as a motivation for a generalization of the  $\beta$ -polynomial concept. Define

$$\beta_t(G, C, x) = \beta_t(G, C) = \alpha(G, x) - 2t\alpha(G - C, x).$$

**Observation 2.** *If  $G$  is a unicyclic graph and  $-1 \leq t \leq 1$ , then  $\beta_t(G, C)$  is a real polynomial.*

This result follows from the fact that it is possible to construct a weighted digraph  $G^t$  whose characteristic polynomial coincides with  $\beta_t(G, C)$ . In order to do this, an edge of  $G$  belonging to the circuit  $C$  is to be exchanged by a pair of oppositely directed arcs, having weights  $e^{i\theta}$  and  $e^{-i\theta}$ . The characteristic polynomial of  $G^t$  obeys the relation [6], [9], [12]

$$(1) \quad \phi(G^t) = \alpha(G) - [e^{i\theta} + e^{-i\theta}]\alpha(G - C) = \alpha(G) - 2t\alpha(G - C),$$

where  $t = \cos \theta$ . For  $\theta$  being a real number the adjacency matrix of  $G^t$  is Hermitean and therefore all the zeros of  $\phi(G^t)$  are real-valued.

Examples show that when the parameter  $t$  is less than  $-1$  or greater than  $+1$ , then the  $\beta_t$ -polynomial of a unicyclic graph may have complex-valued zeros. The simplest such example is provided by the triangle, whose  $\beta_t$ -polynomial is  $x^3 - 3x - 2t$ . It is easy to verify that all the three zeros of this polynomial are real if and only if  $|t| \leq 1$ .

The result formulated here as Observation 2 can easily be extended to graphs possessing several circuits, such that no two circuits share a common edge [6], [9]:

**Observation 3.** *If  $G$  is a graph whose no edge belongs to more than one circuit and  $-1 \leq t \leq 1$ , then for all circuits  $C$  of  $G$ ,  $\beta_t(G, C)$  is a real polynomial.*

In what follows we show that the  $\beta$ -polynomials of some other polycyclic graphs are also real for all values of  $t$ ,  $|t| \leq 1$ .

## 2. THE COMPOUND GRAPHS $G_n$ AND $G_n[]$

The one-vertex graph will be denoted by  $K_1$  and its vertex labeled by  $u_0$ .

Let  $H$  be an  $n$ -vertex graph and let  $S$  be a subset of its vertex set. Then by  $H[S]$  we denote the  $(n + 1)$ -vertex graph obtained by connecting all elements of  $S$  with the vertex  $u_0$  of  $K_1$ .

Let  $G$  be a graph and  $v$  and  $w$  its two (not necessarily distinct) vertices. Construct the graph  $G_n$  by taking  $n$  copies ( $n > 1$ ) of  $G$  and joining the vertex  $v$

**Fig. 1**

of the  $(j + 1)$ -th copy to the vertex  $w$  of the  $j$ -th copy,  $j = 1, 2, \dots, n - 1$  and, in addition, the vertex  $v$  of the first copy to the vertex  $w$  of the  $n$ -th copy (see Fig. 1).

The compound graph  $G_n$  has the following obvious property: The deletion from  $G_n$  of any vertex labeled by  $v$  results in the same subgraph; this subgraph will be denoted by  $G_n - v$ .

If  $G$  is a tree then  $G_n$  is unicyclic and its unique circuit will be denoted by  $C^*$ . Since all the vertices of  $G_n$ , labeled by  $v$ , necessarily belong to the circuit  $C^*$ , the subgraph  $G_n - v$  is acyclic (but needs not be connected).

Denote by  ${}_n$  and  ${}_n$  the sets of vertices of  $G_n$  labeled by  $v$  and  $w$ , respectively (see Fig. 1). The cardinalities of these sets are, of course, equal to  $n$ .

**Lemma 1.** For any  $\subseteq {}_n$ ,

$$(2) \quad \alpha(G_n \setminus \setminus) = x\alpha(G_n) - |\setminus|\alpha(G_n - v)$$

where  $|\setminus|$  stands for the cardinality of the set  $\setminus$ .

If  $\setminus = \emptyset$  then  $G_n \setminus \setminus$  is isomorphic to the disconnected graph  $G_n \cup K_1$ . Because of the identity [2], [5]

$$(3) \quad \alpha(H_a \cup H_b) = \alpha(H_a)\alpha(H_b)$$

and the fact that  $\alpha(K_1) = x$ , in the case when  $\setminus$  is the empty set, eq. (1) is satisfied in a trivial manner.

It remains, therefore, to examine only the case when the set  $\setminus$  is non-empty.

If  $\epsilon$  is an edge of the graph  $H$ , connecting the vertices  $p$  and  $q$ , then [2], [5],

$$(4) \quad \alpha(H) = \alpha(H - \epsilon) - \alpha(H - p - q).$$

Applying the recurrence relation (4) to an edge of  $G_n \setminus \setminus$ , connecting a vertex from  $\setminus$  with  $u_0$ , we obtain

$$\alpha(G_n \setminus \setminus) = \alpha(G_n \setminus \setminus) - \alpha(G_n - v)$$

where  $\setminus' = \setminus \setminus \{v\}$ . If  $\setminus'$  is non-empty, then the application of (4) has to be repeated to the edge of the graph  $G_n \setminus \setminus'$ , connecting  $u_0$  and  $v'$ ,  $v' \in \setminus'$ . This yields

$$\alpha(G_n \setminus \setminus) = \alpha(G_n \setminus \setminus'') - 2\alpha(G_n - v)$$

where  $\setminus'' = \setminus \setminus \{v, v'\}$ . If  $\setminus''$  is non-empty then the procedure has to be repeated again, etc. Ultimately we arrive at

$$\alpha(G_n \setminus \setminus) = \alpha(G_n \cup K_1) - |\setminus|\alpha(G_n - v).$$

Lemma 1 follows now from eq. (3).  $\square$

In a fully analogous manner we may deduce

**Lemma 2.** *Let  $G - v$  and  $G - w$  be isomorphic graphs and  $v \neq w$ . Then for any  $\subseteq_n$  and  $\subseteq_n$*

$$\alpha(G_n[\cup]) = x\alpha(G_n) - (|| + ||)\alpha(G_n - v).$$

According to Lemmas 1 and 2 the matching polynomials of  $G_n[]$  and  $G_n[\cup]$  are independent of the actual choice of the vertices to which the vertex  $u_0$  is connected and depend only on their number, i.e. on the cardinalities of the sets and .

### 3. AN AUXILIARY WEIGHTED DIGRAPH

Denote by  $\phi(H) = \phi(H, x)$  the characteristic polynomial of the graph  $H$ . Suppose that the vertices  $p$  and  $q$  of  $H$  are connected by an edge  $e$  and that the weight of this edge is  $k$ . A well known result of graph spectral theory [3] is the relation

$$(5) \quad \phi(H) = \phi(H - e) - k^2\phi(H - p - q),$$

which holds provided  $e$  is a bridge. (Recall that an edge  $e$  is said to be a bridge of the graph  $H$  if  $H - e$  has more components than  $H$ .) Another well known identity for the characteristic polynomial is [3]

$$(6) \quad \phi(H_a \cup H_b) = \phi(H_a)\phi(H_b).$$

Let  $^1$  be a one-element subset of  $_n$ . Then  $G_n[^1]^k$  denotes the graph obtained from  $G_n[^1]$  by associating the weight  $k$  to the edge  $e_0$  which connects the vertex  $u_0$  with the vertex  $v \in ^1$ . Observe that  $e_0$  is a bridge.

As already pointed out, if  $G$  is a tree then  $G_n$  is unicyclic and its unique circuit is denoted by  $C^*$ . If  $G$  is a tree then  $G_n^t[^1]^k$  denotes the digraph obtained from  $G_n[^1]^k$  by exchanging an edge (any edge) belonging to  $C^*$  by a pair of oppositely directed arcs, having weights  $e^{i\theta}$  and  $e^{-i\theta}$ ,  $t = \cos \theta$ .

Applying eq. (5) to the edge  $e_0$  of  $G_n^t[^1]^k$  and bearing in mind eq. (6) as well as  $\phi(K_1) = x$ , we immediately arrive at

**Lemma 3.** *If  $G$  is a tree then*

$$\phi(G_n^t[^1]^k) = x\phi(G_n^t) - k^2\phi(G_n - v).$$

Further, as a special case of eq. (1) we have

**Lemma 4.** *If  $G$  is a tree then*

$$\phi(G_n^t) = \alpha(G_n) - 2t\alpha(G_n - C^*).$$

#### 4. THE MAIN RESULTS

**Theorem 1.** *If  $G$  is a tree, then for all  $t$ ,  $-1 \leq t \leq 1$ , all  $n > 1$  and all  $\subseteq_n$ ,*

$$\beta_t(G_n \square, C^*) = \phi(G_n^t [1]^k)$$

with  $k = \sqrt{\|\cdot\|}$ .

**Theorem 2.** *Let  $G - v$  and  $G - w$  be isomorphic graphs and  $v \neq w$ . If  $G$  is a tree, then for all  $t$ ,  $-1 \leq t \leq 1$ , all  $n > 1$ , all  $\subseteq_n$  and all  $\subseteq_n$ ,*

$$\beta_t(G_n \sqcup, C^*) = \phi(G_n^t [1]^k)$$

with  $k = \sqrt{\|\cdot\| + \|\cdot\|}$ .

From the definition of the  $\beta_t$ -polynomial,

$$\beta_t(G_n \square, C^*) = \alpha(G_n \square) - 2t\alpha(G_n - C^* \cup K_1).$$

Using eq. (3) and the fact that  $\alpha(K_1) = x$  we get

$$\beta_t(G_n \square, C^*) = \alpha(G_n \square) - 2xt\alpha(G_n - C^*)$$

which combined with Lemma 1 yields

$$(7) \quad \beta_t(G_n \square, C^*) = x[\alpha(G_n) - 2t\alpha(G_n - C^*)] - \|\alpha(G_n - v).$$

If a graph  $H$  is acyclic, then [2], [5],  $\phi(H) \equiv \alpha(H)$ . Consequently,  $\alpha(G_n - v) \equiv \phi(G_n - v)$ . Bearing this fact in mind and using Lemma 4, the right-hand side of (7) is readily transformed into

$$\beta_t(G_n \square, C^*) = x\phi(G_n^t) - \|\alpha(G_n - v).$$

Theorem 1 follows now from Lemma 3.  $\square$

Proof of Theorem 2 is analogous, except that instead of Lemma 1 we now have to employ Lemma 2.  $\square$

All the zeros of the characteristic polynomial of the auxiliary weighted digraph  $G_n^t [1]^k$  are real-valued [3]. Therefore we have

**Corollary 1.1.** *Under the conditions specified in Theorem 1,  $\beta_t(G_n \square)$  is a real polynomial.*

**Corollary 2.1.** *Under the conditions specified in Theorem 2,  $\beta_t(G_n \sqcup)$  is a real polynomial.*

Theorems 1 and 2 can be further extended. Let the graphs  $G_n \square T$  and  $G_n^t [1]^k T$  be obtained by identifying the vertex  $u_0$  of  $G_n \square$  and  $G_n^t [1]^k$ , respectively, with the root of a rooted tree  $T$ . Then we can prove the following results.

**Theorem 3.** *Let  $T$  be an arbitrary rooted tree. Theorem 1 remains valid if  $G_n \square$  and  $G_n^t [1]^k$  are exchanged by  $G_n \square T$  and  $G_n^t [1]^k T$ , respectively.*

**Theorem 4.** *Let  $T$  be an arbitrary rooted tree. Theorem 2 remains valid if  $G_n[\cup]$  and  $G_n^t[1]^k$  are exchanged by  $G_n[\cup]$  and  $G_n^t[1]^k T$ , respectively.*

Theorems 3 and 4 imply that under conditions specified in Theorems 1 and 2,  $\beta_t(G_n[T, C^*])$  and  $\beta_t(G_n[\cup]T, C^*)$  are real polynomials for all rooted trees  $T$  and for all values of the parameter  $t$ ,  $-1 \leq t \leq 1$ .

**Remark.** The special case of Theorems 1 and 3, when the vertices  $v$  and  $w$  coincide and when  $t = 1$  was previously reported by the author in [7].

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Faculty of Science,  
University of Kragujevac,  
P.O. Box 60, 34000 Kragujevac,  
Yugoslavia

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